

# Linked Lists

# Linked List Basics

- Linked lists and arrays are similar since they both store collections of data.
- The *array's* features all follow from its strategy of allocating the memory for all its elements in one block of memory.
- *Linked lists* use an entirely different strategy: linked lists allocate memory for each element separately and only when necessary.

# Disadvantages of Arrays

## 1. The size of the array is fixed.

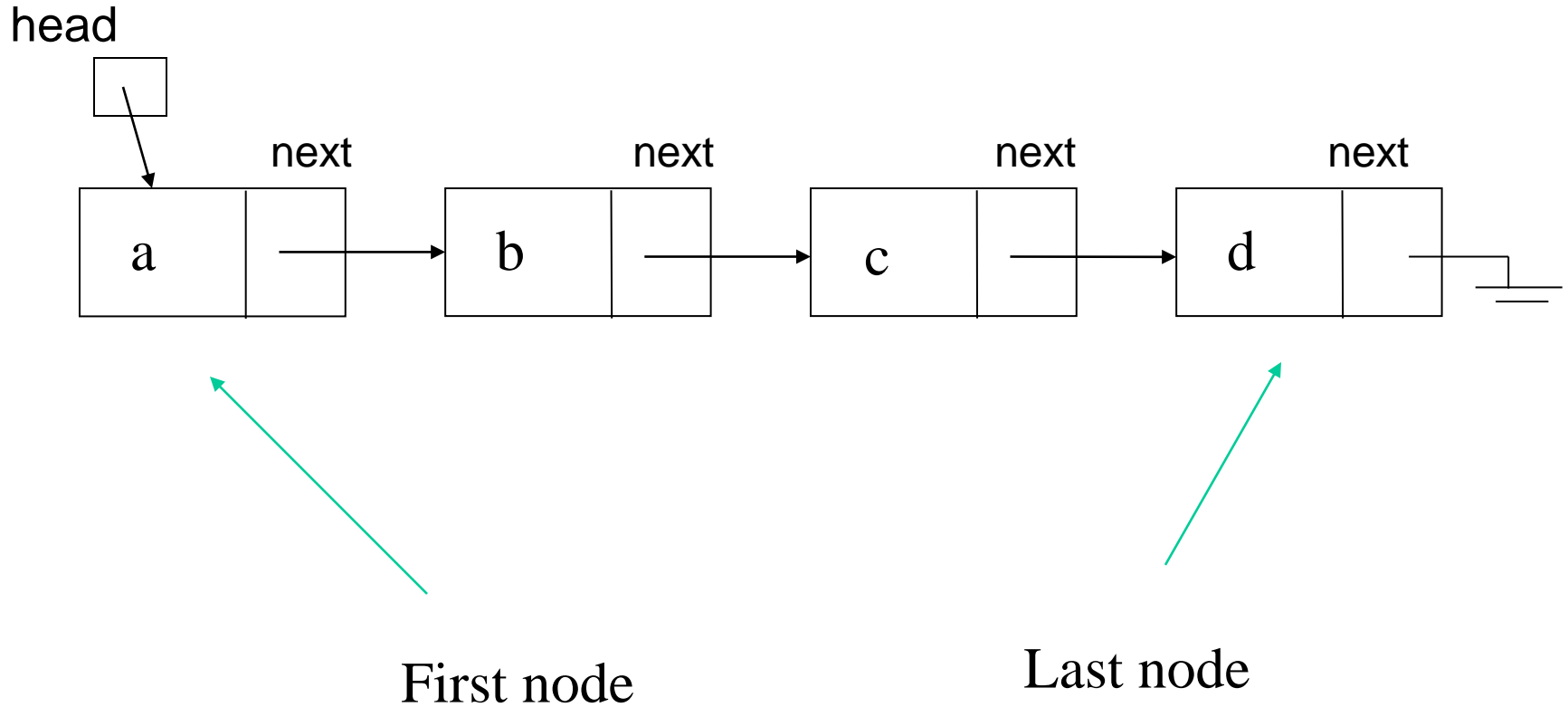
- In case of **dynamically resizing** the array from size  $S$  to  $2S$ , we need  $3S$  units of available memory.
- Programmers allocate arrays which seem "**large enough**" This strategy has two disadvantages: (a) most of the time there are just 20% or 30% elements in the array and 70% of the space in the array really is wasted. (b) If the program ever needs to process more than the declared size, the code breaks.

## 2. Inserting (and deleting) elements into the middle of the array is potentially expensive because existing elements need to be shifted over to make room

# Linked lists

- Linked lists are appropriate when the number of data elements to be represented in the data structure at once is unpredictable.
- Linked lists are dynamic, so the length of a list can increase or decrease as necessary.
- Each node does not necessarily follow the previous one physically in the memory.
- Linked lists can be maintained in sorted order by inserting or deleting an element at the proper point in the list.

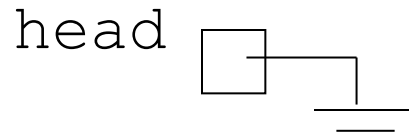
# Singly Linked Lists



# Empty List

- Empty Linked list is a single pointer having the value of NULL.

```
head = NULL;
```



# Basic Ideas

- Let's assume that the node is given by the following type declaration:

```
struct Node {  
    Object element;  
    Node *next;  
};
```

# Basic Linked List Operations

- List Traversal
- Searching a node
- Insert a node
- Delete a node

# Traversing a linked list

```
Node *pWalker;
int count = 0;

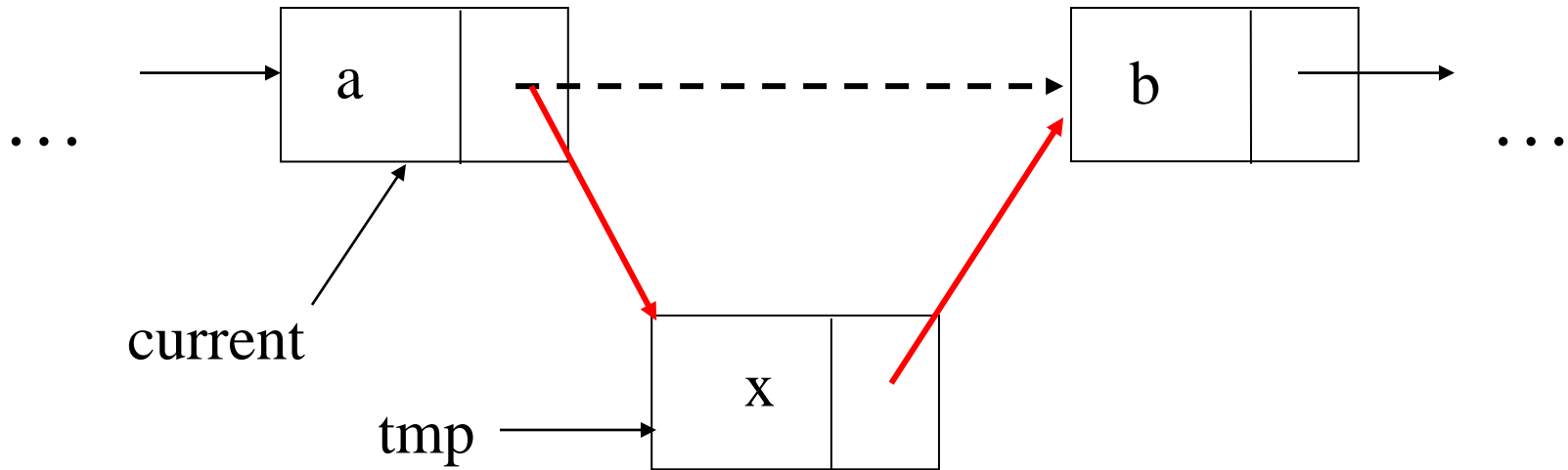
cout << "List contains:\n";

for (pWalker=pHead; pWalker!=NULL;
     pWalker = pWalker->next)
{
    count ++;
    cout << pWalker->element << endl;
}
```

# Searching a node in a linked list

```
pCur = pHead;  
  
// Search until target is found or we reach  
// the end of list  
while (pCur != NULL &&  
       pCur->element != target)  
{  
    pCur = pCur->next;  
}  
  
//Determine if target is found  
if (pCur) found = 1;  
else found = 0;
```

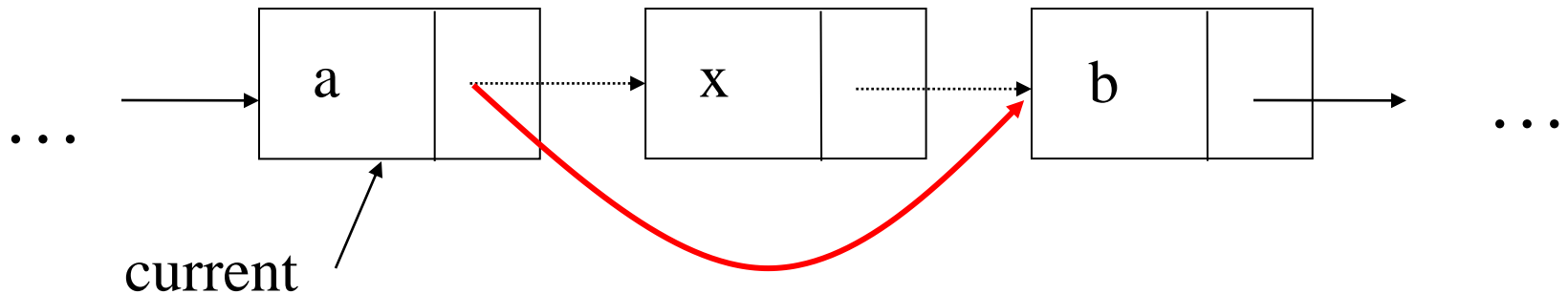
# Insertion in a linked list



```
tmp = new Node;  
tmp->element = x;  
tmp->next = current->next;  
current->next = tmp;
```

Or simply (if Node has a constructor initializing its members):  
`current->next = new Node(x, current->next);`

# Deletion from a linked list



```
Node *deletedNode = current->next;  
current->next = current->next->next;  
delete deletedNode;
```

# Special Cases (1)

- Inserting before the first node (or to an empty list):

```
tmp = new Node;
tmp->element = x;
if (current == NULL) {
    tmp->next = head;
    head = tmp;
}
else { // Adding in middle or at end
    tmp->next = curent->next;
    current->next = tmp;
}
```

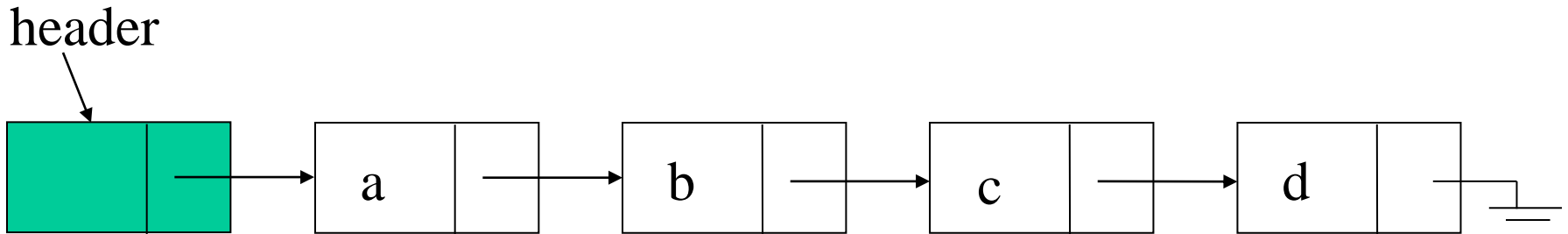
# Special Cases (2)

```
Node *deletedNode;
if (current == NULL) {
    // Deleting first node
    deletedNode = head;
    head = head ->next;
}
else{
    // Deleting other nodes
    deletedNode = current->next;
    current->next = deletedNode ->next;
}
delete deletedNode;
```

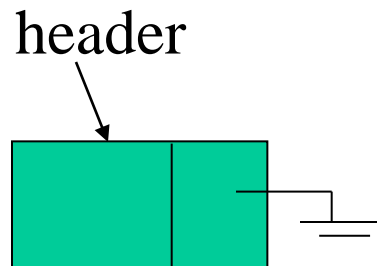
# Header Nodes

- One problem with the basic description: it assumes that whenever an item  $x$  is removed (or inserted) some previous item is always present.
- Consequently removal of the first item and inserting an item as a new first node become special cases to consider.
- In order to avoid dealing with special cases: introduce a **header node (dummy node)**.
- A header node is an extra node in the list that holds no data but serves to satisfy the requirement that every node has a previous node.

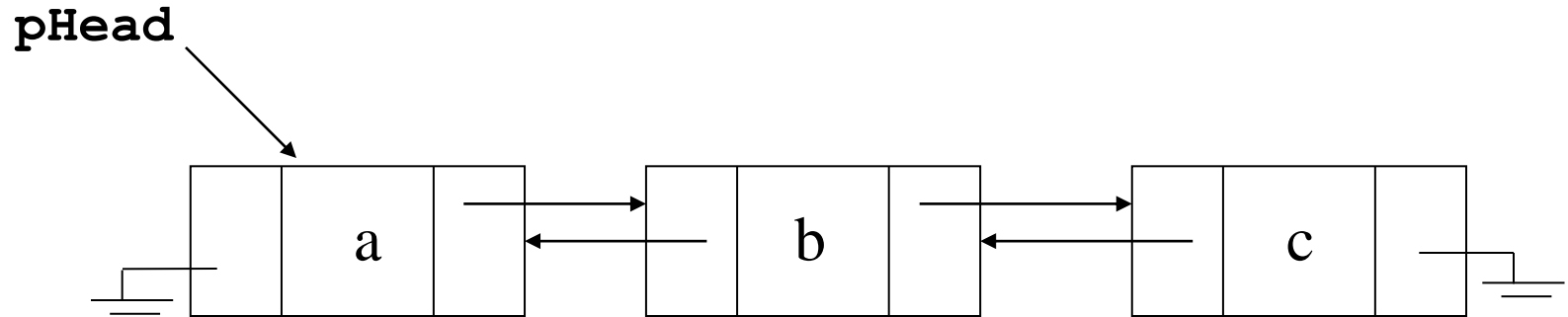
# List with a header node



## Empty List



# Doubly Linked Lists



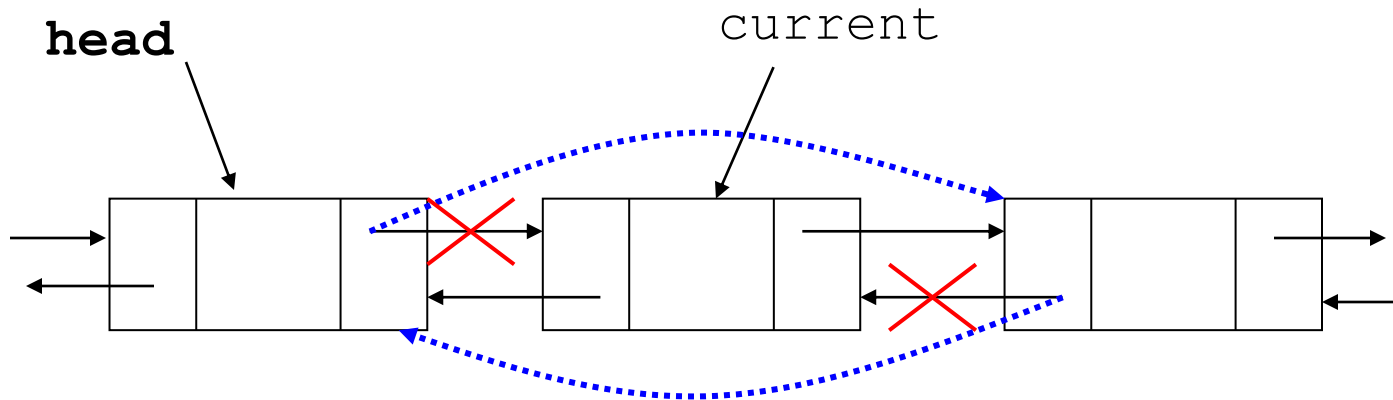
## Advantages:

- Convenient to traverse the list backwards.
- Simplifies insertion and deletion because you no longer have to refer to the previous node.

## Disadvantage:

- Increase in space requirements.

# Deletion

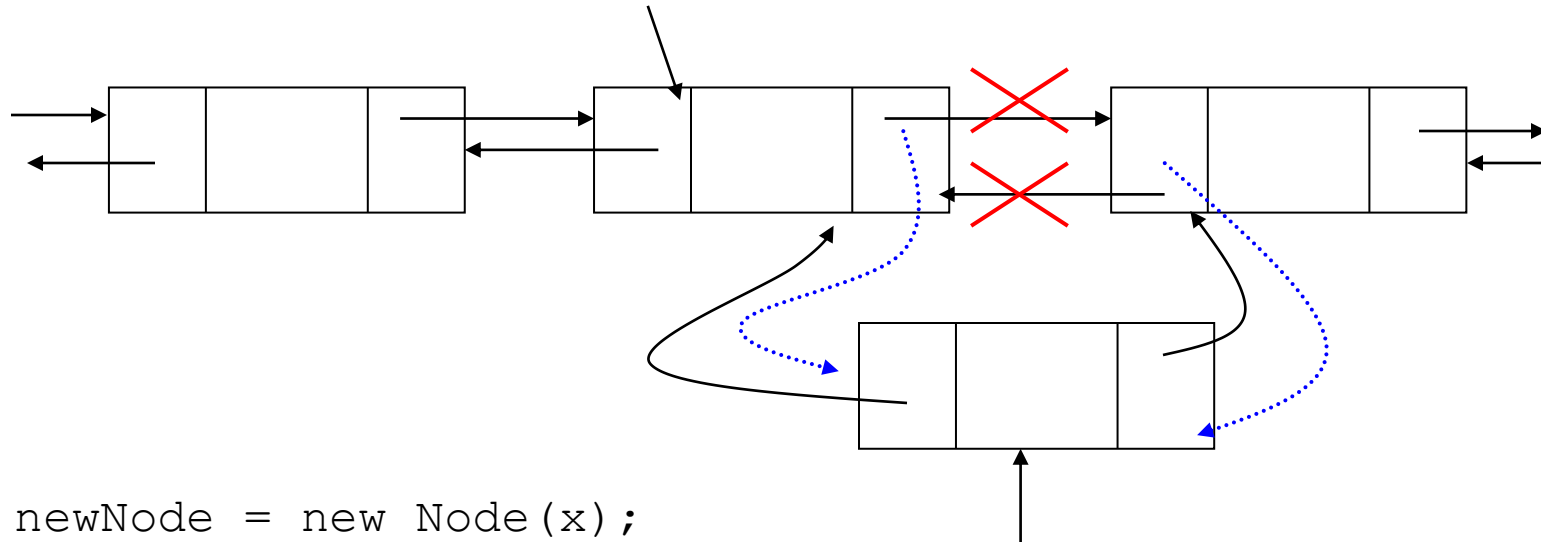


```
oldNode = current;  
oldNode->prev->next = oldNode->next;  
oldNode->next->prev = oldNode->prev;  
delete oldNode;  
current = head;
```

# Insertion

head

current

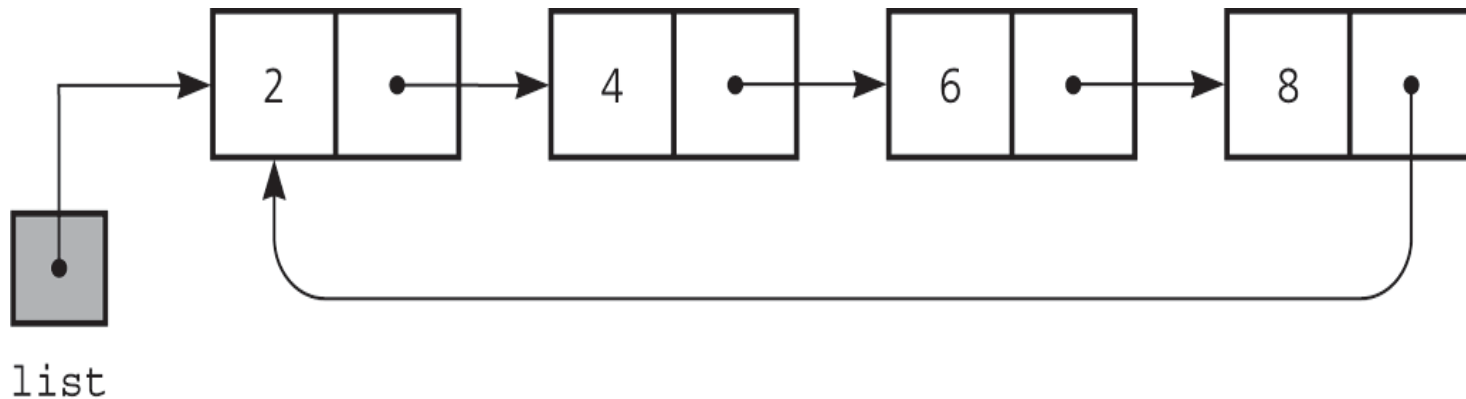


newNode

```
newNode = new Node(x);  
newNode->prev = current;  
newNode->next = current->next;  
newNode->prev->next = newNode;  
newNode->next->prev = newNode;  
current = newNode;
```

# Circular Linked Lists

- Last node references the first node
- Every node has a successor
- No node in a circular linked list contains *NULL*



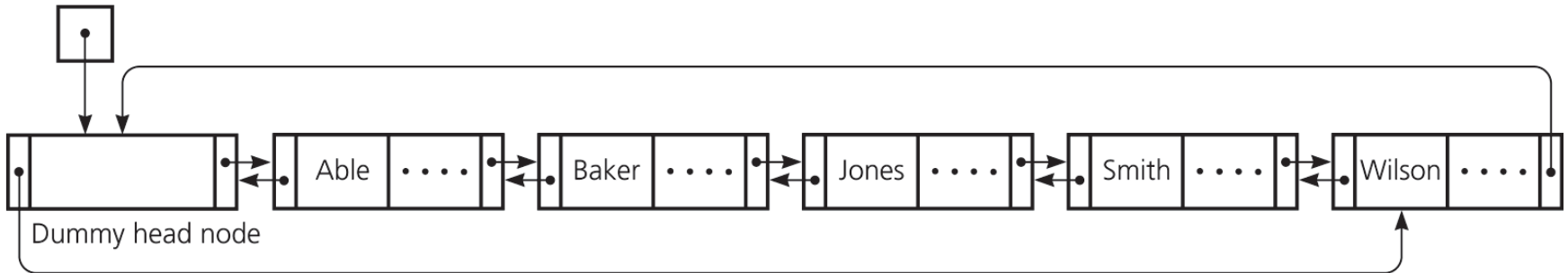
A circular linked list

# Circular Doubly Linked Lists

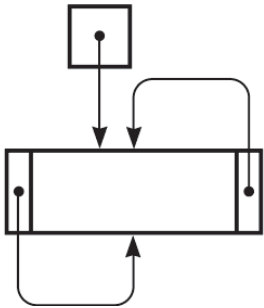
- Circular doubly linked list
  - `prev` pointer of the dummy head node points to the last node
  - `next` reference of the last node points to the dummy head node
  - No special cases for insertions and deletions

# Circular Doubly Linked Lists

(a) listHead



(b) listHead



(a) A circular doubly linked list with a dummy head node

(b) An empty list with a dummy head node